# Quantum Foundations Lecture 14

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HSC112

#### Announcements

- Schmid College Academic Advising:
  - Tuesday April 3, 4:30pm-6:30pm AF209A (Prof. Leifer)
  - Wednesday April 4, 4:30pm-6:30pm Henley Hall Basement (Prof. Dressel)
- Adam Becker is returning to Chapman:
  - Book event and signing at 1888 center: Monday April 16. RSVP required <a href="https://bit.ly/AdamBecker">https://bit.ly/AdamBecker</a>
- Assignments
  - First Draft due on Blackboard April 11.
  - Peer review until April 16.
  - Discussion in class April 16.
  - Final Version due May 2.
- Homework 3 due April 11.
- I like lunch invitations

# Review of Last Lecture

# The View from the Smaller Space

# Trace Preservation

# Trace Preservation

# Positivity vs. Complete Positivity

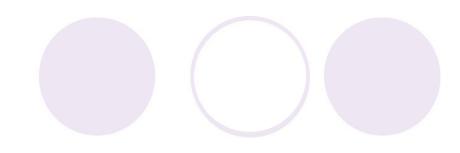
# Positivity vs. Complete Positivity

# Complete Positivity

# Complete Positivity

# Complete Positivity

# Summary



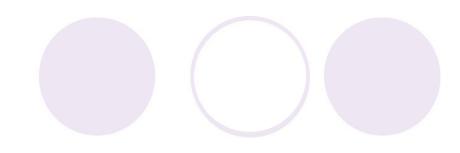
# Examples of Qubit CPT Maps

# 8.v) Postivie Operator Valued Measures (POVMs)

# The View from the Larger Church

# The View from the Larger Church

# Summary



# Example



# Application: Minimum Error Discrimination

- Alice has a preparation device that prepares the system in either the state  $\rho$  or the state  $\sigma$ . She chooses each with 50/50 probability and sends the system to Bob.
- $\bullet$  Bob makes a measurement on the system and has to guess whether  $\rho$  or  $\sigma$  was prepared.
- What is his maximum probability of success and what measurement should he make?

- Let's look at the classical case first. There is a variable that can take d possible values  $j=1,2,\cdots,d$ .
- $\odot$  Alice prepares the probability distributions p or q with 50/50 probability.
- Bob decides on a subset  $E_p \subseteq \{1,2,\cdots,d\}$ . If  $j \in E_p$  he guesses p. If it is in the complement  $E_q = \{1,2,\cdots,d\} \setminus E_p$ , he guesses q.

- What is his probability of success:
  - $p_{\text{SUCC}} = \text{Prob}(\boldsymbol{p} \text{ is prepared}) \text{Prob}(E_{\boldsymbol{p}}|\boldsymbol{p}) + \text{Prob}(\boldsymbol{q} \text{ is prepared}) \text{Prob}(E_{\boldsymbol{q}}|\boldsymbol{q})$   $= \frac{1}{2} \left[ \sum_{j \in E_n} p_j + \sum_{j \in E_n} q_j \right]$
- ullet However,  $\sum_{j \in E_{m{q}}} q_j = 1 \sum_{k \in E_{m{p}}} q_j$  and so

$$p_{\text{succ}} = \frac{1}{2} \left[ 1 + \sum_{j \in E_p} (p_j - q_j) \right]$$

Therefore,

$$p_{\text{succ}} \le \frac{1}{2} \left[ 1 + \sum_{\{j \mid p_j \ge q_j\}} (p_j - q_j) \right]$$

which can be achieved if Bob guesses  $\boldsymbol{p}$  whenever  $p_j \geq q_j$  and  $\boldsymbol{q}$  otherwise.

We can rewrite this in a simpler way by noting that

$$\textstyle \sum_{\{j \mid p_j \geq q_j\}} p_j + \sum_{\{j \mid p_j < q_j\}} p_j = 1 \quad \text{and} \quad \sum_{\{j \mid p_j \geq q_j\}} q_j + \sum_{\{j \mid p_j < q_j\}} q_j = 1$$

Subtracting these and rearranging gives

$$\sum_{\{j|p_j \ge q_j\}} (p_j - q_j) = \sum_{\{j|p_j < q_j\}} (q_j - p_j)$$

• Therefore,

$$p_{\text{succ}} \leq \frac{1}{2} \left[ 1 + \sum_{\{j \mid p_j \geq q_j\}} (p_j - q_j) \right]$$

$$= \frac{1}{2} \left[ 1 + \frac{1}{2} \left( \sum_{\{j \mid p_j \geq q_j\}} (p_j - q_j) + \sum_{\{j \mid p_j < q_j\}} (q_j - p_j) \right) \right]$$

$$= \frac{1}{2} \left[ 1 + \frac{1}{2} \sum_{j=1}^{d} |p_j - q_j| \right]$$

$$= \frac{1}{2} [1 + D_c(\mathbf{p}, \mathbf{q})]$$

• Where  $D_c(\boldsymbol{p},\boldsymbol{q}) = \frac{1}{2}\sum_{j=1}^d \left| p_j - q_j \right|$  is called the variational distance.

### Quantum Case

 Theorem (Helstrom, Holevo): The optimal success probability in the quantum case is given by

$$p_{\text{succ}} = \frac{1}{2} \left[ 1 + D_q(\rho, \sigma) \right]$$

where

$$D_q(\rho, \sigma) = \frac{1}{2} \text{Tr}(|\rho - \sigma|)$$

is known as the trace distance, and the matrix norm is given by

$$|M| = \sqrt{M^{\dagger}M}$$
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