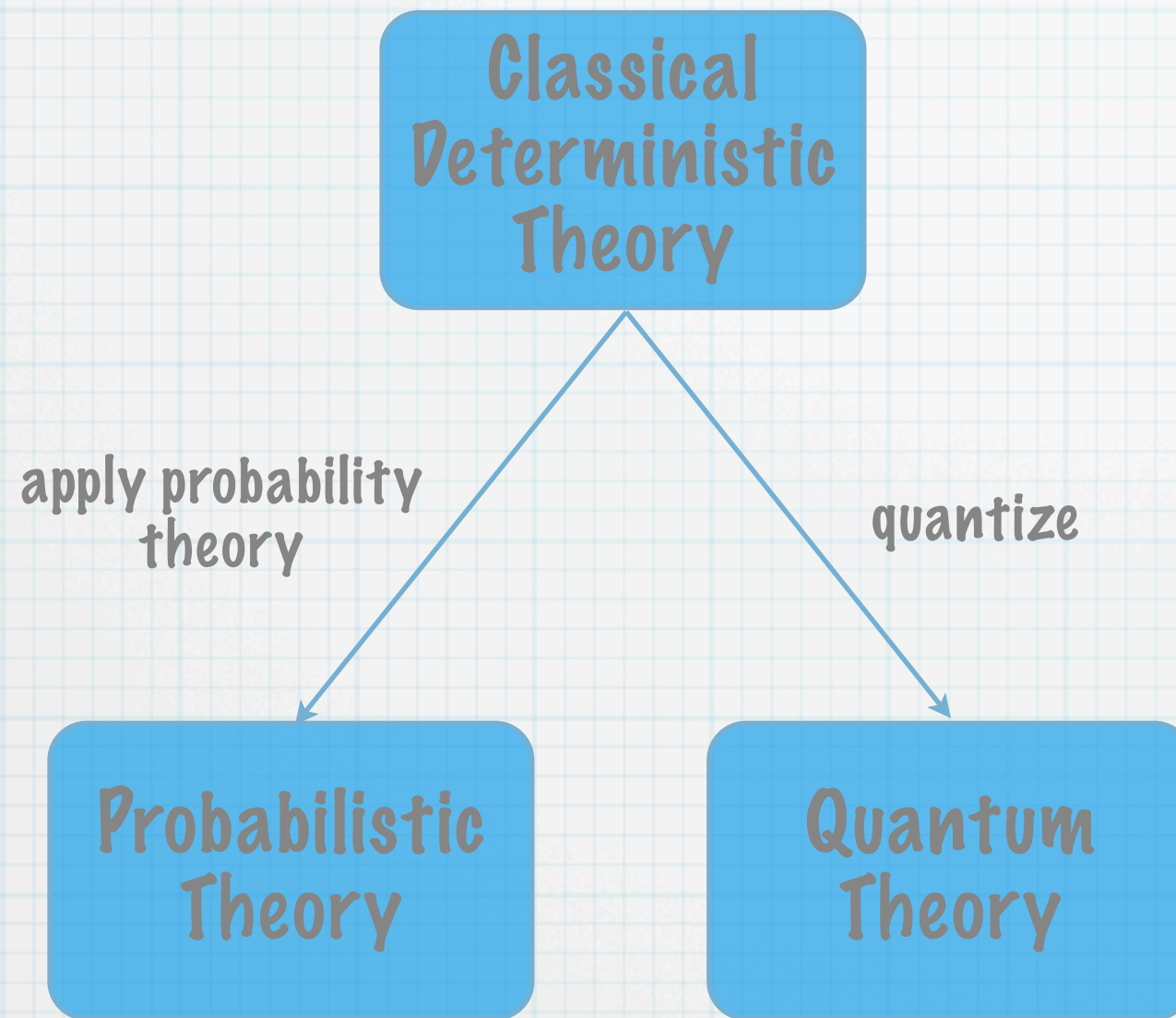


Quantum Dynamics as Generalized Conditional Probabilities

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FPP-4 Växjö (8th June 2006)



Quantum Theory as a Meta-theory



- * If it's that simple, why is it hard to quantize GR?
- * Applies to quantum measurement theory, but something is missing.
- * QT not as abstract as PT. Causal structure is still present in QT.

Classical Probability vs. Quantum Theory

Classical	Quantum
Probability distribution: $P(X)$	Quantum State: ρ_A
Joint distribution: $P(X, Y)$	Joint State: ρ_{AB}
Transition matrix: $\Gamma_{Y X}$	TPCP map: $\mathcal{E}_{B A}$
Conditional Prob.: $P(Y X)$?

Why quantum conditional probability?

- * Conditional probabilities allow all types of correlation to be treated on an equal footing, whether timelike, spacelike or completely abstract.
 - * Causal relations are not primitive in probability theory.
- * Some classical probabilistic structures are defined in terms of conditional probability.
 - * Markov Chains
 - * Bayesian Networks
- * Some Bayesians take conditional probability to be the most fundamental notion.
 - * See textbook by D. V. Lindley

Outline

1. Introduction

- i. The Many faces of conditional probability
- ii. Suggestions for a quantum analog of conditional probability

2. Stochastic Dynamics as Conditional Probabilities

3. Choi-Jamiołkowski Isomorphism

4. A New Isomorphism

5. Operational Interpretation

6. Application: Cloning, broadcasting & monogamy of entanglement

7. Future Directions

1. Introduction

(A) Reconstructing a joint distribution from a marginal

$$P(X, Y) = P(Y|X)P(X)$$

(B) Bayesian Updating

$$P(H|D) = \frac{P(D|H)P(H)}{P(D)}$$

(C) Stochastic Dynamics

$$P(Y = i) = \sum_j (\Gamma_{Y|X})_{ij} P(X = j)$$

(D) Conditional Shannon Entropy

$$H(Y|X) = - \sum_{X,Y} P(X, Y) \log_2 P(Y|X)$$

(E) Reduction of complexity via conditional independence

$$P(Y|X, Z) = P(Y|Z) \Leftrightarrow P(X, Y, Z) = P(X|Z)P(Y|Z)P(Z)$$

1. Introduction

(A) Reconstruction of a joint state ρ_{AB} from a marginal ρ_A .

(B) Updating quantum states after a measurement

$$\rho|_M = \frac{\mathcal{E}^M(\rho)}{\text{Tr}(M\rho)}$$

(C) TPCP dynamics

$$\rho_B = \mathcal{E}_{B|A}(\rho_A)$$

(D) Conditional von Neumann Entropy

$$S(B|A) = -\text{Tr}(\rho_{AB} \log_2 \rho_{B|A})$$

(E) Reduction of complexity via conditional independence

1. Introduction

* Cerf & Adami ('97-'99):

$$\begin{aligned}\rho_{B|A} &= 2^{\log_2 \rho_{AB} - \log_2 \rho_A \otimes I_B} \\ &= \lim_{n \rightarrow \infty} \left[\rho_{AB}^{\frac{1}{n}} (\rho_A \otimes I_B)^{\frac{1}{n}} \right]^n\end{aligned}$$

* (A) Reconstruction:

$$\rho_{AB} = 2^{\log_2 \rho_A \otimes I_B + \log_2 \rho_{B|A}}$$

* (C) Entropy:

$$S(B|A) = S(A, B) - S(A) = -\text{Tr} (\rho_{AB} \log_2 \rho_{B|A})$$

* (E) Complexity Reduction: If $\log_2 \rho_{B|AC} = I_A \otimes \log_2 \rho_{B|C}$

$$\rho_{ABC} = 2^{I_{AB} \otimes \log_2 \rho_C + \log_2 \rho_{A|C} \otimes I_B + I_A \otimes \log_2 \rho_{B|C}}$$

1. Introduction

- * (B) Updating:

- * POVM: $M = \{M\}, \quad M > 0, \quad \sum_M M = I$

- * Probability Rule: $P(M) = \text{Tr}(M\rho)$

- * Update CP-map: $\rho|_M = \frac{\mathcal{E}^M(\rho)}{\text{Tr}(M\rho)}$

$$\mathcal{E}^M(\rho) = \sum_j A_j^M \rho A_j^{M\dagger} \quad \sum_j A_j^{M\dagger} A_j^M = M$$

- * \mathcal{E}^M depends on details of system-measuring device interaction.

1. Introduction

- * Is there one update rule that is more “Bayes’ rule like” than the rest?

- * Traditionally (see Bub ‘77 for projective measurements):

$$\rho|_M = \frac{\sqrt{M} \rho \sqrt{M}}{\text{Tr}(M \rho)}$$

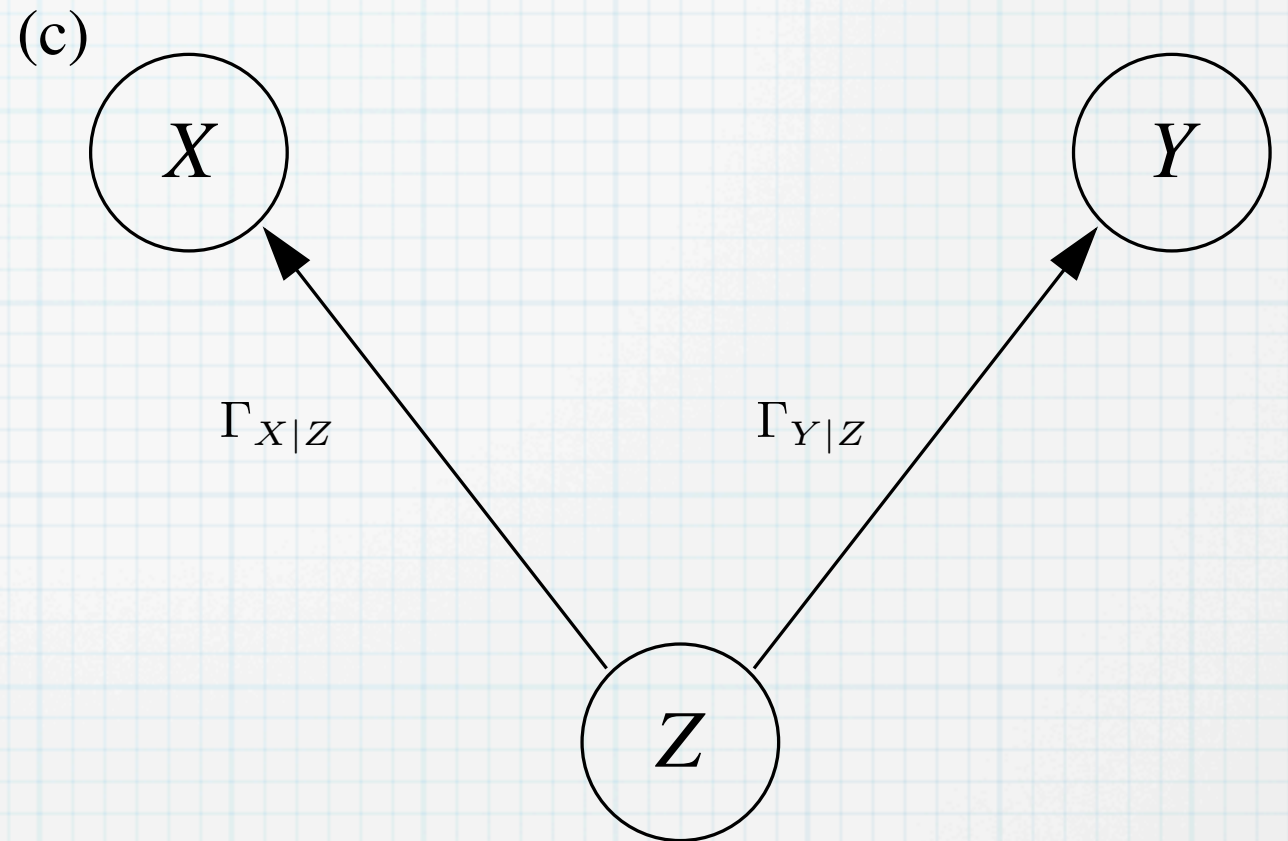
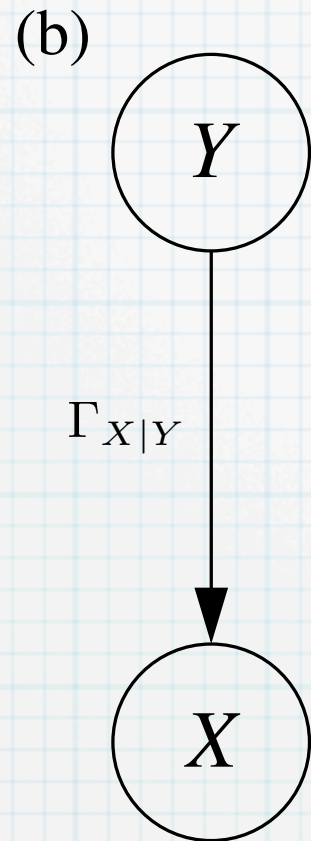
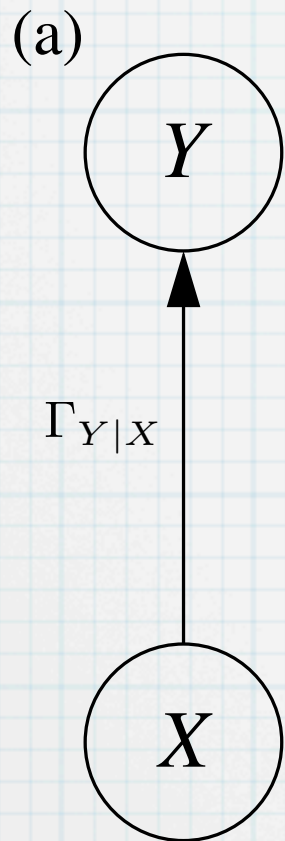
- * According to Fuchs (‘01, ‘02):

$$\rho|_M = \frac{\sqrt{\rho} M \sqrt{\rho}}{\text{Tr}(M \rho)}$$

- * Both reduce to Bayes’ rule when the M are projection operators and

$$[M, \rho] = 0$$

2. Dynamics as conditional probability



$$P(Y) = \sum_X \Gamma_{Y|X} P(X)$$

$$P(X, Y) = \Gamma_{Y|X} P(X)$$

$$P(Y|X) = \Gamma_{Y|X}^r$$

$$P(X, Y)$$

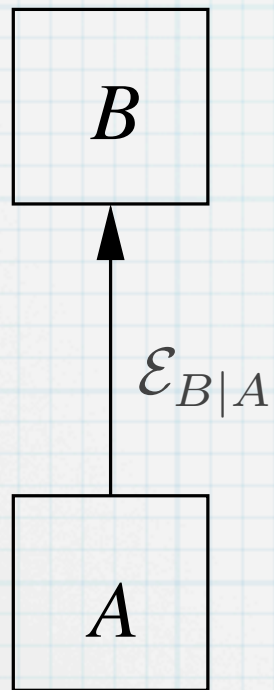
$$P(X) = \sum_Y P(X, Y)$$

$$P(Y|X) = \frac{P(X, Y)}{P(X)}$$

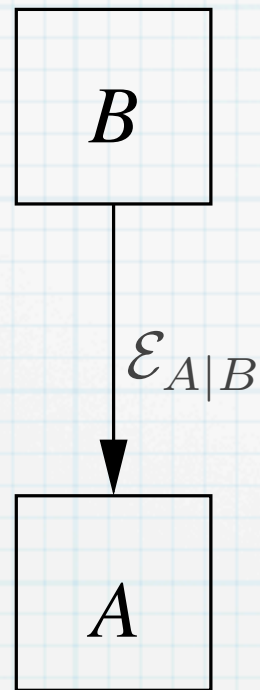
Isomorphism: $(P(X), \Gamma_{Y|X}^r) \Leftrightarrow P(X, Y)$

2. Dynamics as conditional probability

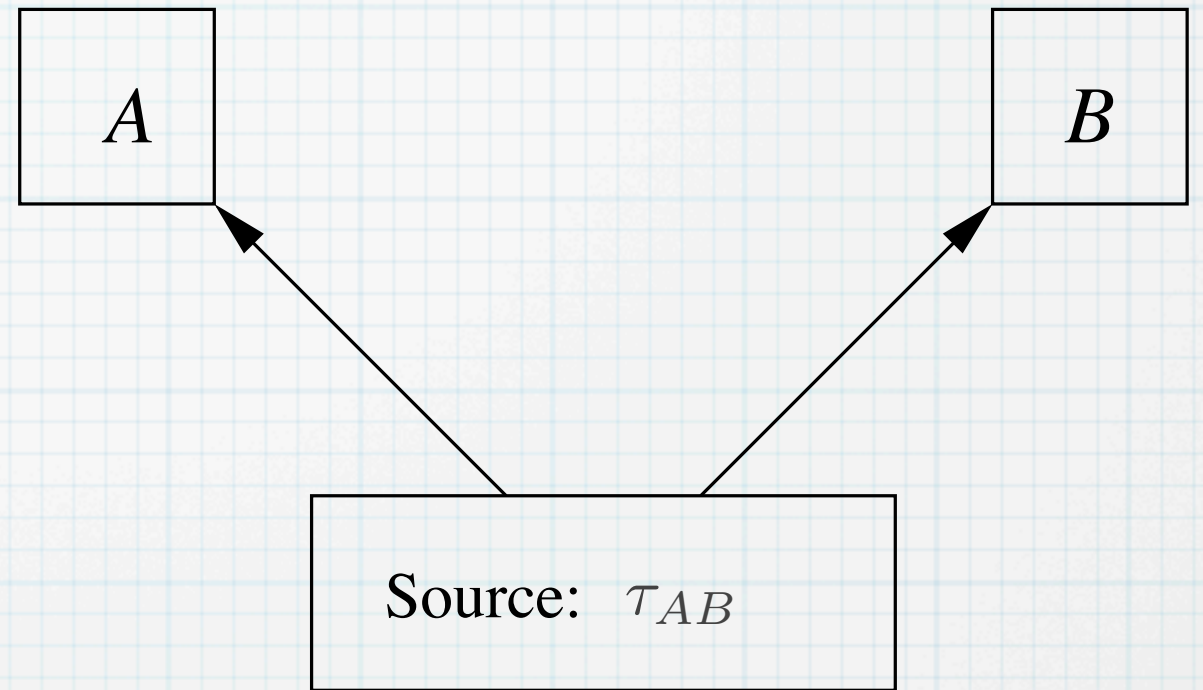
(a)



(b)



(c)



$$\rho_B = \mathcal{E}_{B|A}(\rho_A)$$

$$\rho_{AB} = ?$$

$$\rho_{B|A}, \mathcal{E}_{B|A}^r, ?$$

$$\tau_{AB}$$

$$\tau_A = \text{Tr}_B(\tau_{AB})$$

$$\tau_{B|A} = ?$$

$$\text{Isomorphism: } \left(\rho_A, \mathcal{E}_{B|A}^r \right) \Leftrightarrow \tau_{AB} \text{ ?}$$

3. Choi-Jamiołkowski Isomorphism

* For bipartite pure states and operators:

$$R_{B|A} = \sum_{jk} \alpha_{jk} |j\rangle_B \langle k|_A \Leftrightarrow |\Psi\rangle_{AB} = \sum_{jk} \alpha_{jk} |k\rangle_A \otimes |j\rangle_B$$

* For mixed states and CP-maps:

$$\mathcal{E}_{B|A}(\rho_A) = \sum_{\mu} R_{B|A}^{(\mu)} \rho_A R_{B|A}^{(\mu)\dagger} \Rightarrow \tau_{AB} = \sum_{\mu} \left| \Psi^{(\mu)} \right\rangle_{AB} \left\langle \Psi^{(\mu)} \right|_{AB}$$

3. Choi-Jamiołkowski Isomorphism

- * Let $|\Phi^+\rangle_{AA'} = \frac{1}{\sqrt{d_A}} \sum_j |j\rangle_A \otimes |j\rangle_{A'}$

- * Then $\tau_{AB} = \mathcal{I}_A \otimes \mathcal{E}_{B|A'} (|\Phi^+\rangle_{AA'} \langle \Phi^+|_{AA'})$

$$\mathcal{E}_{B|A}(\rho_A) = d_A^2 \langle \Phi^+|_{AA'} \rho_A \otimes \tau_{A'B} |\Phi^+\rangle_{AA'}$$

- * Operational interpretation: Noisy gate teleportation.

3. Choi-Jamiołkowski Isomorphism

* Remarks:

- * Isomorphism is basis dependent. A basis must be chosen to define $|\Phi^+\rangle_{AA'}$.
- * If we restrict attention to Trace Preserving CP-maps then

$$\tau_A = \text{Tr}_B (\tau_{AB}) = \frac{1}{d_A} I_A$$

- * This is a special case of the isomorphism we want to construct

$$\left(\rho_A, \mathcal{E}_{B|A}^r \right) \Leftrightarrow \tau_{AB}$$

where $\rho_A = \frac{1}{d_A} I_A$.

4. A New Isomorphism

* $\left(\rho_A, \mathcal{E}_{B|A}^r\right) \rightarrow \tau_{AB}$ direction:

* Instead of $|\Phi^+\rangle_{AA'}$, use $|\Phi\rangle_{AA'} = (\rho_A^T)^{\frac{1}{2}} \otimes I_{A'} |\Phi^+\rangle_{AA'}$

* Then $\tau_{AB} = \mathcal{I}_A \otimes \mathcal{E}_{B|A}^r (|\Phi\rangle_{AA'} \langle\Phi|_{AA'})$

* $\tau_{AB} \rightarrow \left(\rho_A, \mathcal{E}_{B|A}^r\right)$ direction:

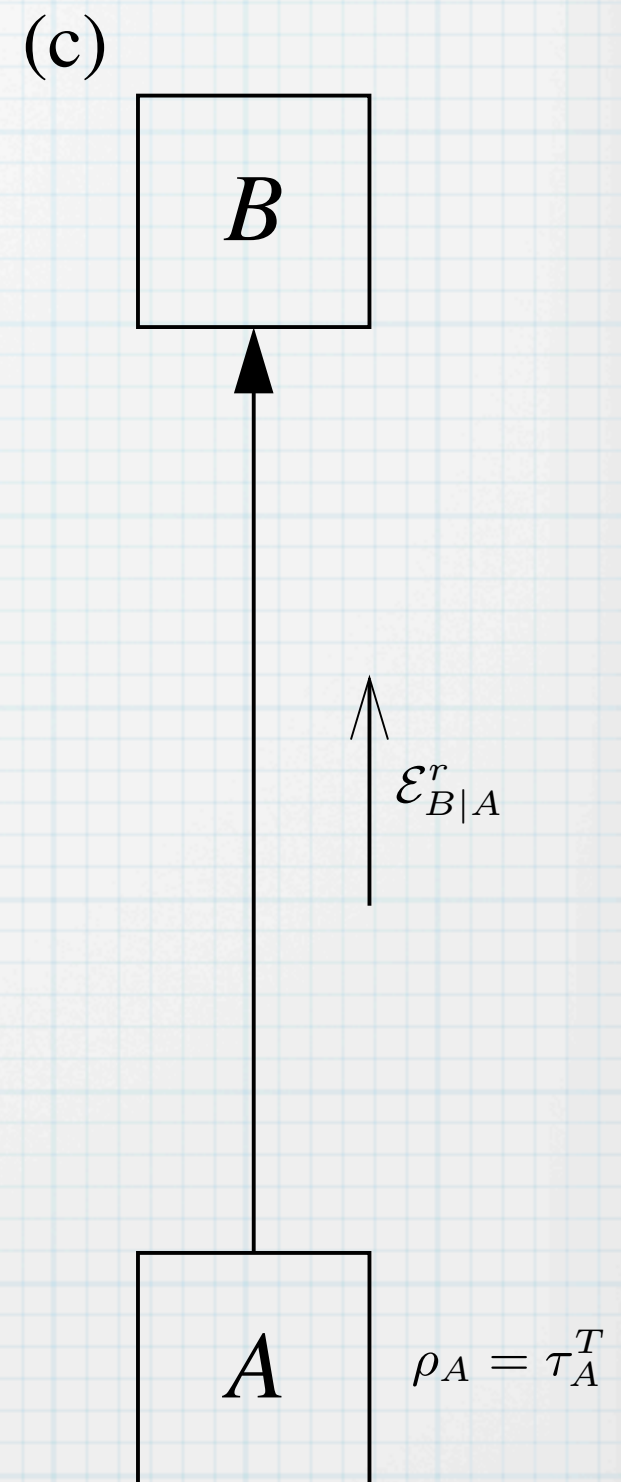
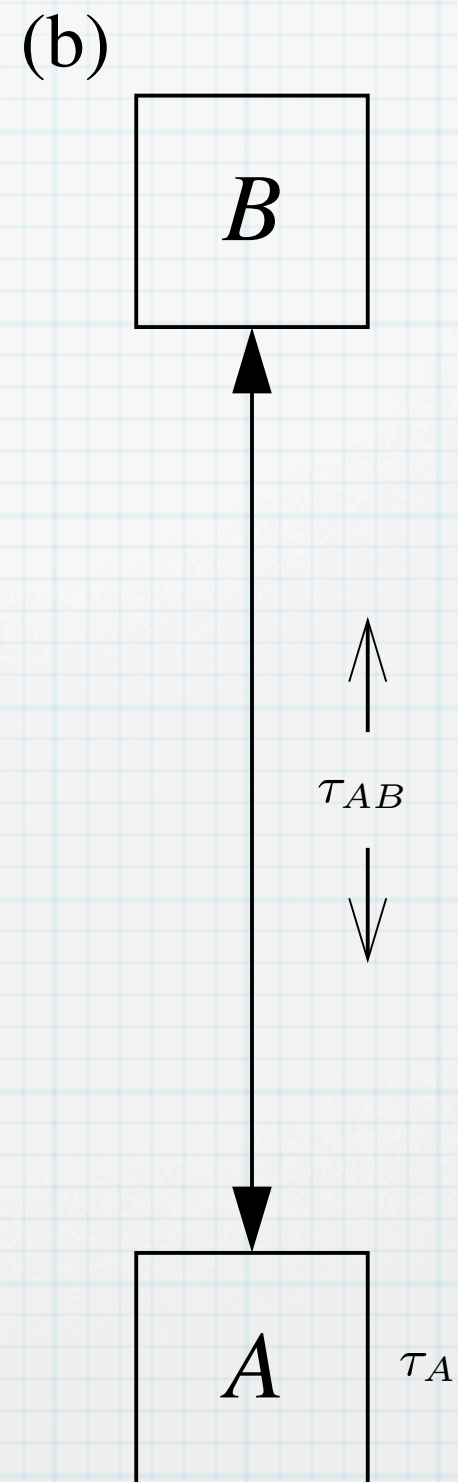
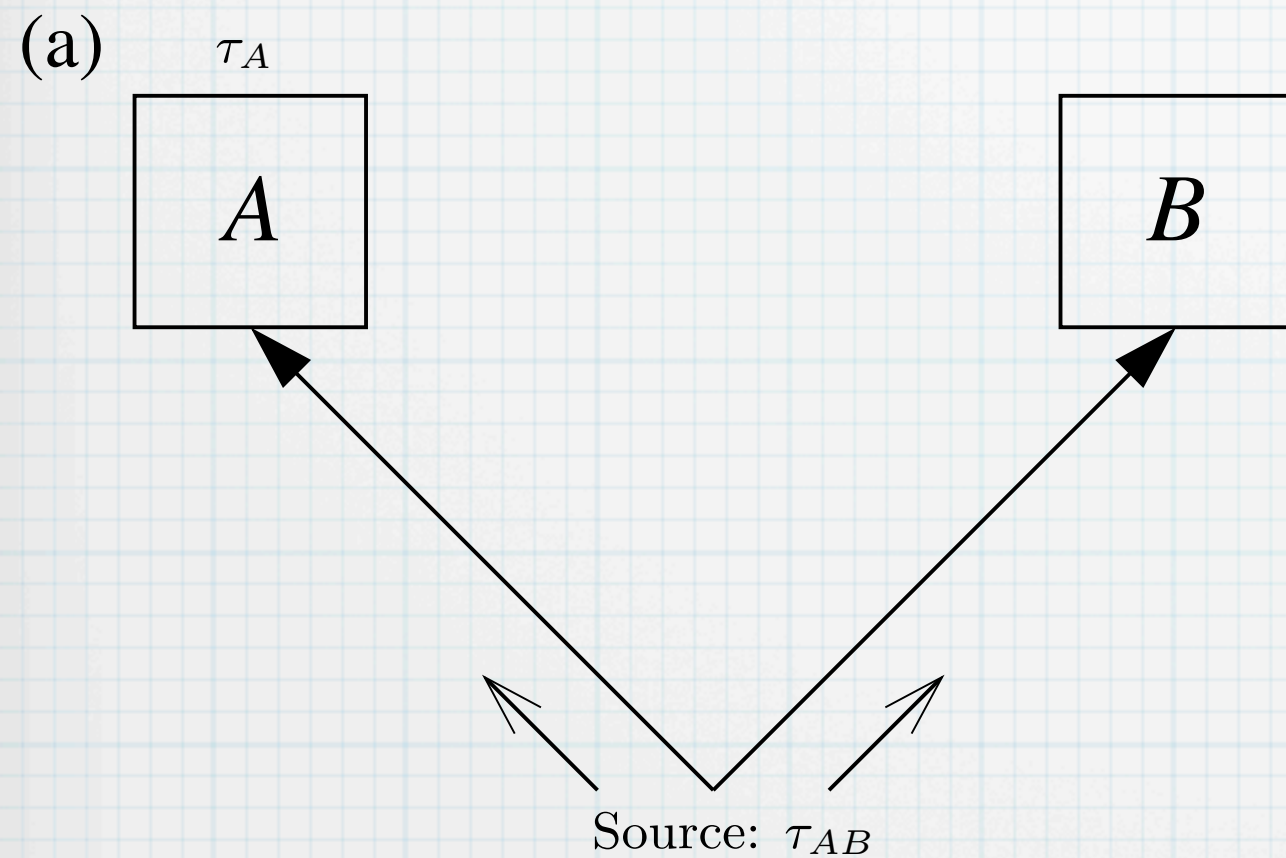
* Set $\rho_A = \tau_A^T$, $\tau_A = \text{Tr}_B (\tau_{AB})$

* let $\sigma_{B|A} = \tau_A^{-\frac{1}{2}} \otimes I_B \tau_{AB} \tau_A^{-\frac{1}{2}} \otimes I_B$

* $\sigma_{B|A}$ is a density operator, satisfying $\text{Tr}_B (\sigma_{B|A}) = \frac{1}{d_A^r} P_A$

* It is uniquely associated to a TPCP map $\mathcal{E}_{B|A}^r : \mathcal{L}(P_A \mathcal{H}_A) \rightarrow \mathcal{L}(\mathcal{H}_B)$
via the Choi-Jamiołkowski isomorphism.

4. A New Isomorphism



5. Operational Interpretation

- * Reminder about measurements:

- * POVM: $M = \{M\}, \quad M > 0, \quad \sum_M M = I$

- * Probability Rule: $P(M) = \text{Tr}(M\rho)$

- * Update CP-map: $\rho|_M = \frac{\mathcal{E}^M(\rho)}{\text{Tr}(M\rho)}$

$$\mathcal{E}^M(\rho) = \sum_j A_j^M \rho A_j^{M\dagger} \quad \sum_j A_j^{M\dagger} A_j^M = M$$

- * \mathcal{E}^M depends on details of system-measuring device interaction.

5. Operational Interpretation

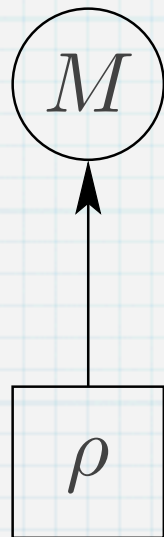
* Lemma: $\rho = \sum_M P(M) \rho_{|M}$ is an ensemble decomposition of a

density matrix ρ iff there is a POVM $M = \{M\}$ s.t.

$$P(M) = \text{Tr}(M\rho) \qquad \rho_{|M} = \frac{\sqrt{\rho} M \sqrt{\rho}}{\text{Tr}(M\rho)}$$

* Proof sketch: $M = P(M) \rho^{-\frac{1}{2}} \rho_{|M} \rho^{-\frac{1}{2}}$

5. Operational Interpretation



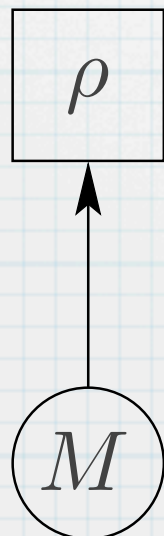
* M -measurement of ρ

* Input: ρ

* Measurement probabilities: $P(M) = \text{Tr}(M\rho)$

* Updated state:

$$\rho|_M = \frac{\sqrt{M}\rho\sqrt{M}}{\text{Tr}(M\rho)}$$



* M -preparation of ρ

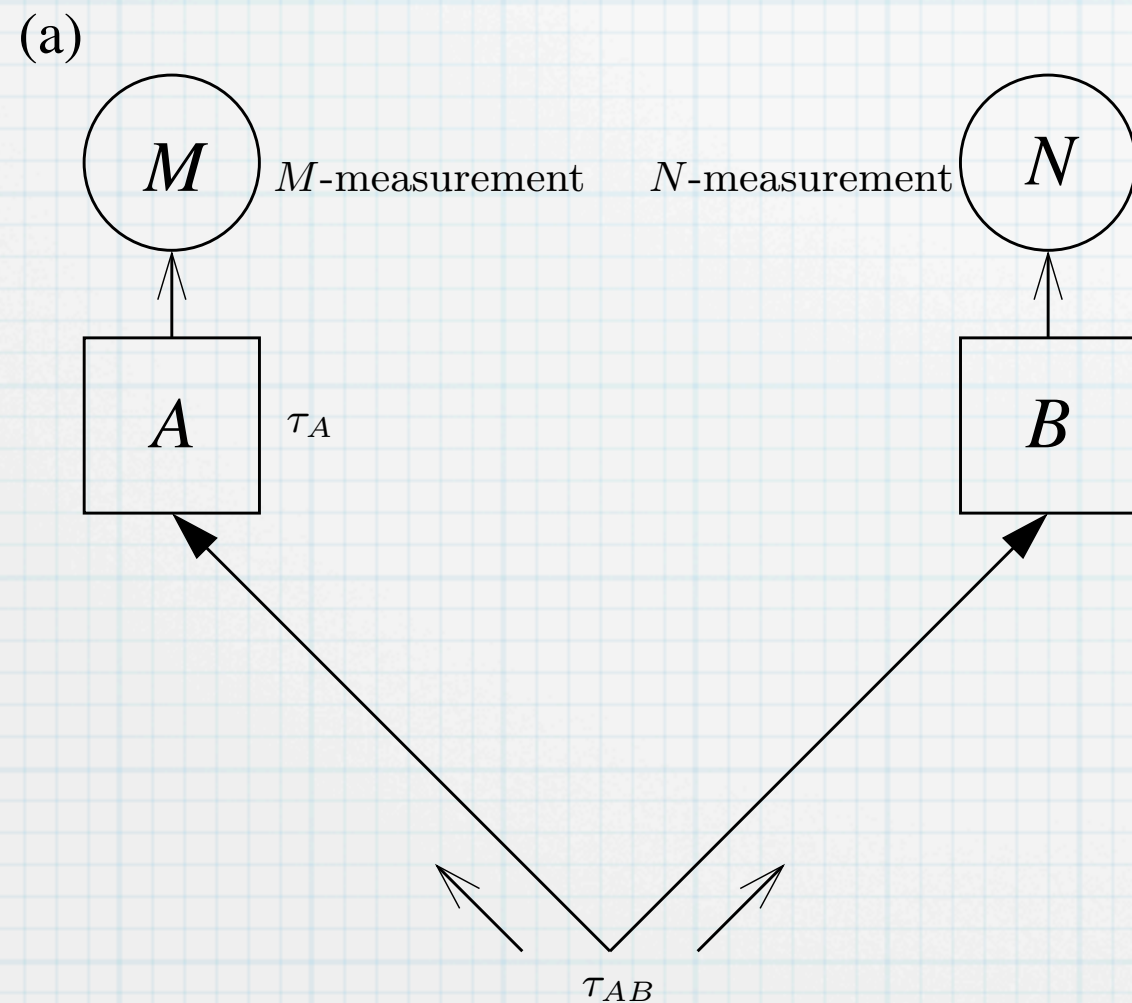
* Input: Generate a classical r.v. with p.d.f

$$P(M) = \text{Tr}(M\rho)$$

* Prepare the corresponding state:

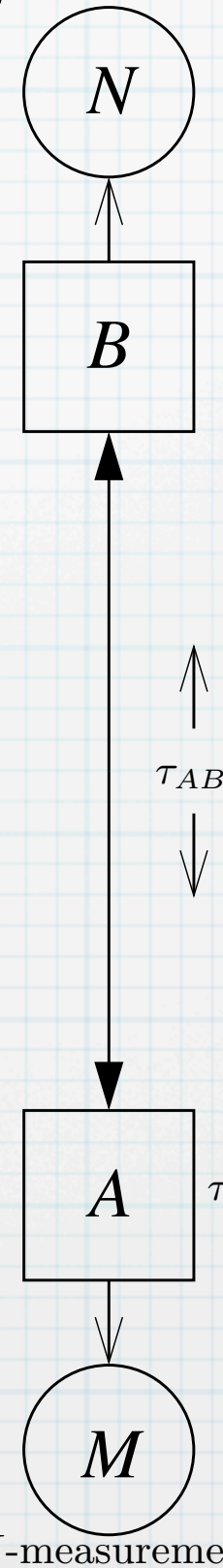
$$\rho|_M = \frac{\sqrt{\rho}M\sqrt{\rho}}{\text{Tr}(M\rho)}$$

5. Operational Interpretation

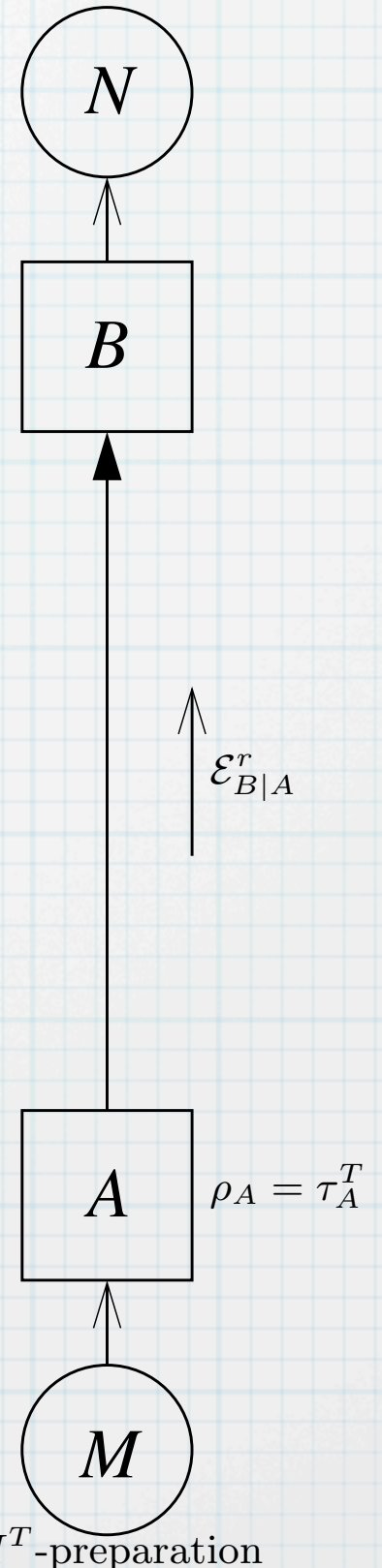


$P(M, N)$ is the same in (a) and (c) for any POVMs M and N .

(b) N -measurement



(c) N -measurement



5. Application: Broadcasting & Monogamy

- * For any TPCP map $\mathcal{E}_{BC|A} : \mathcal{L}(\mathcal{H}_A) \rightarrow \mathcal{L}(\mathcal{H}_B \otimes \mathcal{H}_C)$ the reduced maps are:

$$\mathcal{E}_{B|A} = \text{Tr}_C \circ \mathcal{E}_{BC|A} \quad \mathcal{E}_{C|A} = \text{Tr}_B \circ \mathcal{E}_{BC|A}$$

- * The following commutativity properties hold:

$$\begin{array}{ccc} \rho_{ABC} & \xlongequal{\quad} & (\rho_A, \mathcal{E}_{BC|A}^r) \\ \text{Tr}_C \downarrow & & \downarrow \text{Tr}_C \\ \rho_{AB} & \xlongequal{\quad} & (\rho_A, \mathcal{E}_{B|A}^r). \end{array}$$

- * Therefore, 2 states ρ_{AB}, ρ_{AC} incompatible with being the reduced states of a global state ρ_{ABC} .
- * 2 reduced maps $\mathcal{E}_{B|A}^r, \mathcal{E}_{C|A}^r$ incompatible with being the reduced maps of a global map $\mathcal{E}_{BC|A}^r$.

5. Application: Broadcasting & Monogamy

- * A TPCP-map $\mathcal{E}_{A'A''|A} : \mathcal{L}(\mathcal{H}_A) \rightarrow \mathcal{L}(\mathcal{H}_{A'} \otimes \mathcal{H}_{A''})$ is broadcasting for a state ρ_A if

$$\mathcal{E}_{A'|A}(\rho_A) = \rho_{A'} \quad \mathcal{E}_{A''|A}(\rho_A) = \rho_{A''}$$

- * A TPCP-map $\mathcal{E}_{A'A''|A} : \mathcal{L}(\mathcal{H}_A) \rightarrow \mathcal{L}(\mathcal{H}_{A'} \otimes \mathcal{H}_{A''})$ is cloning for a state ρ_A if

$$\mathcal{E}_{A'A''|A}(\rho_A) = \rho_{A'} \otimes \rho_{A''}$$

- * Note: For pure states cloning = broadcasting.
- * A TPCP-map is universal broadcasting if it is broadcasting for every state.

5. Application: Broadcasting & Monogamy

- * No cloning theorem (Dieks '82, Wootters & Zurek '82):
 - * There is no map that is cloning for two nonorthogonal and nonidentical pure states.
- * No broadcasting theorem (Barnum et. al. '96):
 - * There is no map that is broadcasting for two noncommuting density operators.
- * Clearly, this implies no universal broadcasting as well.
- * Note that the maps $\mathcal{E}_{A'|A}$, $\mathcal{E}_{A''|A}$ are valid individually, but they cannot be the reduced maps of a global map $\mathcal{E}_{A'A''|A}$.

5. Application: Broadcasting & Monogamy

- * The maps $\mathcal{E}_{A'|A}, \mathcal{E}_{A''|A}$ must be related to incompatible states $\tau_{AA'}, \tau_{AA''}$
- * Theorem: If $\mathcal{E}_{A'A''|A}$ is universal broadcasting, then both $\tau_{AA'}, \tau_{AA''}$ must be pure and maximally entangled.
- * Ensemble broadcasting $\{(p, \rho_1), ((1-p), \rho_2)\}$ s.t. $[\rho_1, \rho_2] \neq 0$
$$\left(p\rho_1 + (1-p)\rho_2, \mathcal{E}_{A'A''|A}^r\right) \Leftrightarrow \tau_{AA'A''}$$
- * Theorem: There is a local operation on A that transforms both $\tau_{AA'}$ and $\tau_{AA''}$ into pure, entangled states with nonzero probability of success.

7. Future Directions

- * Quantitative relations between approximate ensemble broadcasting and monogamy inequalities for entanglement.
- * More generally, useful in analyzing any qinfo protocol involving the action of a TPCP-map on a particular ensemble rather than the whole Hilbert space.
- * Can the various analogs of conditional probability be unified?
- * Can quantum theory be developed using an analog of conditional probability as the fundamental notion?
- * Can we eliminate background causal structures entirely from the formalism of quantum theory?